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Project	IEEE 802.16 Broadband Wireless Access Working Group <http: 16="" ieee802.org=""></http:>		
Title	Routing with CID Encapsulation		
Date Submitted	2007-03-05		
Source(s)	Aik Chindapol Jimmy Chui Hui Zeng Siemens Corporate Research Princeton, NJ, USA	Voice: +1 609 734 3364 Fax: +1 609 734 6565 Email: aik.chindapol@siemens.com	
	Byung-Jae Kwak, Sungcheol Chang, Young-il Kim ETRI 161, Gajeong-Dong, Yuseong-Gu, Daejeon, Korea 305-350	Voice: +82-42-860-6618 Fax: +82-42-861-1966 bjkwak@etri.re.kr	
	Kyu Ha Lee, Changkyoon Kim, Hyung Kee Kim Samsung Thales San 14, Nongseo-Dong, Giheung-Gu,	Voice: +82-31-280-9917 Fax: +82-31-280-1562 kyuha.lee@samsung.com	
	Yongin, Gyeonggi-Do, Korea 449-712 Yuan-Ying Hsu Telcordia Applied Research Center Taiwan Co., Taipei, Taiwan	Voice: +886-2-37895177#4558 Fax: +886-2-26552078 yyhsu@tarc-tw.research.telcordia.com	
Re:	This is in response to the call for technical comments and contributions for IEEE 802.16j		
Abstract	This document describes fast connection routing via CID encapsulation.		
Purpose	This contribution is provided as input for the IEEE 802.16j baseline document.		
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Fast Connection Establishment and Maintenance with Relays

2 Introduction

1

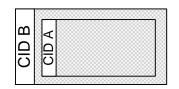
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3 In 802.16e, each connection (both management and data) is identified by a Connection ID (CID).

Connections are identified by a 16-bit CID [2]. At MS initialization, management connections (basic connection
and primary management connection) are established. The basic connection is for short, time-urgent MAC
management messages and the primary management connection is for longer, more delay tolerant MAC
management messages. In addition, transport connections for downlink and uplink are established to transfer
data flows. There is no routing required; data is transmitted solely between the BS and the MS.

9 In a centralized multi-hop relay system, the routing for each MS is decided by the BS. The routing path is 10 based on a number of factors such as measured channel qualities, QoS of each connection, fairness, etc. Each 11 RS must be informed of which packets to detect and forward in order to provide the correct route for the 12 packets. For wireless mobile networks, however, the topology of the network and the channel conditions 13 change rapidly. Therefore, creating the routing structure and maintaining it are quite challenging.

This contribution proposes to allow the encapsulation of CIDs for the relaying of packets, as in [1]. That is, if a packet is to be transmitted to an MS via an RS, the packet intended for the MS (with CID "A") becomes the payload of a second packet (with CID "B") intended for the RS. When the RS receives the data, it understands to only retransmit the embedded data, i.e. the original packet intended for the MS.



19 Figure 1: Example of a CID encapsulated packet

This procedure can be used as a quick method for relaying data. The example in this proposal is the use of CID encapsulation for fast route updating. If the routing path for a particular MS changes, it may be possible to perform a temporary route changes via encapsulation. Although encapsulation incurs additional overhead, it is a quicker method of changing routing path than using other methods due to the time required to transmit and acknowledge management messages.

- 25 The proposal for this implementation has the following advantages:
- 26 Simplified operation of the relay
- 27 Initiating encapsulation operation is time-efficient. It is made by the BS without RS/MS involvement.

28 General Description

In the legacy 802.16e standard, each management and data connection for a particular MS is assigned a connection ID (CID) by the BS. The DL-MAP contains information about the CIDs in each allocated burst. By decoding this DL-MAP, the MS can identify which bursts it should listen to. This functionality should not be changed with the addition of relays.

CID encapsulation is a simple concept where the packet for a target recipient is the payload for an
 intermediary node, as described above. Although RSs do not have data to transmit on its own, the base station

1 assigns a specific CID for each RS. When the intermediary node receives a packet whose CID belongs to itself,

it strips away the outermost-layered header, decodes and forwards the payload with the CID indicated for the
 next hop, which itself is a packet for the next recipient. The allocation of each transmission is allocated by the

4 BS using a regular MAP structure. This procedure can be concatenated to have many outer layers. The number

of layers required is the number of hops (minus one) in the path from BS to MS. This procedure satisfies the

5 of layers required is the number of hops (ninus one) in the path from BS to MS. This procedure satisfies the

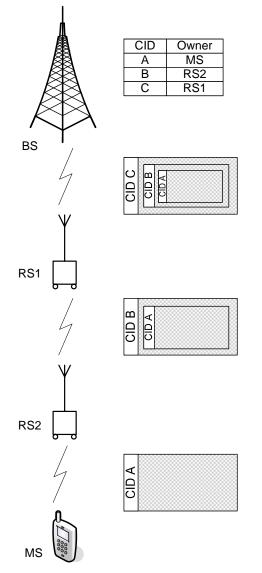
6 requirement of maintaining MS legacy operation.

7 **CID Encapsulation**

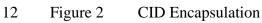
8 For this method, the BS may encapsulate the message with a CID intended for the RS. Once the RS decodes it,

9 it removes encapsulation information and only forwards the original message with a different CID. If the

10 transmission involves multiple hops, the message may have to be encapsulated many times.



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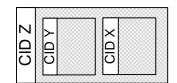
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1 Above is Figure 2 depicting encapsulation. The CIDs for the next hop are embedded in the data payload of the

current hop. There are no requirements for the relationship between the CID of the current packet and the
 embedded CIDs. This allows fast route changes when data transfer is time-critical and there is not enough time

4 to perform proper route update.

5 It is also possible to have more than one embedded packet (multiple packets with multiple CIDs) in the packet 6 containing encapsulated data.



8 **Proposed Text**

9 ------ Beginning of Text Changes ------

10

7

11 6.3.2.1 MAC header formats

12 [Insert the following at the end of 6.3.2.1:]

13 The MAC header of the PDU from the MS to the MR-BS via the RS is encapsulated by the access RS, and the

14 MAC header of the PDU from the MR-BS to the MS via the RS is decapsulated by the access RS.

15

16 [Change the text in Table 4 as indicated:]

17 Table 4 – MAC header format

Syntax	Size	Notes
MAC Header() {		
НТ	1 bit	0 = Generic MAC header
		1 = Bandwidth request header
EC	1 bit	If $HT = 1$, $EC = 0$
if (HT == 0) {		
Туре	6 bits	
СЕ	1 bit	0 = no CID encapsulation
		1 = CID encapsulation is in use
CI	1 bit	
EKS	2 bits	
Reserved	1 bit	Shall be set to zero
LEN	11 bits	
}		

else {		
Туре	3 bits	
BR	19 bits	
}		
CID	16 bits	
HCS	8 bits	
}		

1

2 [Add the following text to 6.3.25]

3 6.3.25 Relay path management and routing

When a relay station receives a MAC PDU with the CE field set in the MAC header, it shall remove the current MAC header and forward the payload as the new PDU. If CRC is used, the BS calculates the CRC for each packet. This reduces the calculation required at each intermediary node.

7 ----- End of Text Changes -----

8

9 **References**

- 10 [1] IEEE C802.16j-06/004r1, "Recommendations on IEEE 802.16j".
- 11 [2] IEEE 802.16-2004, "Part 16: Air Interface for Fixed and Mobile Broadband Wireless Access Systems".
- 12 [3] IEEE 802.16e-2005, "Part 16: Air Interface for Fixed and Mobile Broadband Wireless Access Systems,
- 13 Amendment 2: Physical and Medium Access Control Layers for Combined Fixed and Mobile Operation in
- 14 Licensed Bands and Corrigendum 1".

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