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Title	Service flow management for RS				
Date Submitted	2007-03-05				
Source(s)	Kenji Saito, Takashi Inoue KDDI R&D Laboratories Inc. Hikarino-oka 7-1, Yokosuka, Kanagawa 239-0847, Japan	Voice: +81 46 847 6347 Fax: +81 46 847 0947 saito@kddilabs.jp			
	Shulan Feng, Yanling Lu, Ting Li, Liangliang Zhang Hisilicon Technologies Harbour Building, No.8, Dongbeiwang West Road, HaiDian District, Beijing, China	Voice: 86-10-82829010 Fax: 86-10-82829075 luyanling@hisilicon.com fengsl@huawei.com			
	Yuefeng Zhou Fujitsu Laboratories of Europe Ltd Hayes Park Central, Hayes End Road, Hayes, Middlesex, UB4 8FE, UK	Voice: +44 (0) 20 8573 4444 Fax: +44 (0) 20 8606 4539 Yuefeng.zhou@uk.fujitsu.com			
	Masato Okuda Fujitsu Laboratories Ltd. Kamikodanaka 4-1-1, Kawasaki, 211-8588, Japan	Voice: +81-44-754-2811 Fax: +81-44-754-2786 okuda@jp.fujitsu.com			
Re:	This contribution is a response to "IEEE 802.16j-07/007r2 Contributions regarding IEEE Project 802.16j" (2007-02-19)	Call for Technical Comments and			
Abstract	This document proposes service flow management sequence	e through RS.			
Purpose	This document is provided in response for Call for Technical Co IEEE Project 802.16j	mments and Contributions regarding			
Notice	This document has been prepared to assist IEEE 802.16. It is offered binding on the contributing individual(s) or organization(s). The material in form and content after further study. The contributor(s) reserve(s) material contained herein.	al in this document is subject to change			
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patented technology (or technology under patent application) might be incorporated into a draft standard being developed within the IEEE 802.16 Working Group. The Chair will disclose this notification via the IEEE 802.16 web site http://ieee802.org/16/ipr/patents/notices.

Service flow management for RS

Kenji Saito, Takashi Inoue KDDI R&D Laboratories Inc.

ShuLan Feng, Yanling Lu, Ting Li, Liangliang Zhang Hisilicon Technologies

Yuefeng Zhou
Fujitsu Laboratories of Europe Ltd

Masato Okuda Fujitsu Laboratories Ltd.

1. General

This document presents an amendment for service flow management of RS.

2. Background

In centralized scheduling case, the MRBS manages all subordinate nodes. However, in distributed scheduling case, the MR-BS does not need to manage all nodes, and RS may be able to manage its subordinate node by itself.

Based on the method proposed in [2], the management message sent by the MR-BS can be read and authenticated by the nodes on the mulit-hop link. We propose the service flow management sequence through RS for centralized/distributed scheduling case.

3. Text to be inserted into standard

6.3.14.9.3 DSA

6.3.14.9.3.1 SS-initiated DSA

Insert the following table the end of 6.3.14.9.3.1:

In MR centralized scheduling case, a RS only forwards each DSA messages from SS to MR-BS and vice versa. The MR-BS checks whether the QoS requirements can be supported both on relay link (MR-BS ~ RS) and on access link (RS ~ MS). This process is illustrated in Table 125a.

Table 125a – DSA initiated from SS through RS (Centralized scheduling case)

SS		<u>RS</u>		<u>MR-</u> BS
New service flow needed				
Check if resource are available	le			
Send DSA-REQ	DSA-REQ>	Receive / Send DSA-REQ	DSA-REQ>	Receive DSA-REQ
Set Timers T7 and T14				
Timer T14 Stops	<dsx-rvd< td=""><td>Receive / Send DSX-RVD</td><td><dsx-rvd< td=""><td>DSA-REQ integrity valid</td></dsx-rvd<></td></dsx-rvd<>	Receive / Send DSX-RVD	<dsx-rvd< td=""><td>DSA-REQ integrity valid</td></dsx-rvd<>	DSA-REQ integrity valid
				Check whether SS is authorized for Service
				Check whether service flow QoS can be supported both on relay link and on access link
				Create SFID
				If uplink AdmittedQoSParamSet is non-null, map service flow to CID
				If uplink ActiveQoSParamSet is non-null, Enable reception of data on new uplink service flow
Receive DSA-RSP	<dsa-rsp< td=""><td>Receive / Send DSA-RSP</td><td><dsa-rsp< td=""><td>Send DSA-RSP</td></dsa-rsp<></td></dsa-rsp<>	Receive / Send DSA-RSP	<dsa-rsp< td=""><td>Send DSA-RSP</td></dsa-rsp<>	Send DSA-RSP
Timer T7 Stops				
If ActiveQoSParamSet is no transmission and/or reception new service flow				
Send DSA-ACK	DSA-ACK>	Receive / Send DSA-ACK	DSA-ACK>	Receive DSA-ACK
				If downlink ActiveQoSParamSet is non-null, Enable transmission of data on new downlink service flow

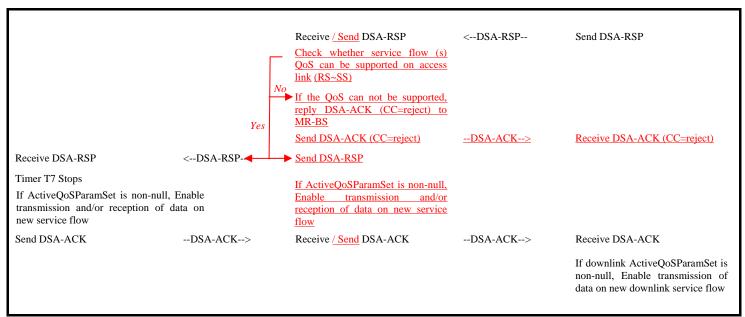
^a Authorization happens prior to the DSA-REQ being received by the BS. The details of BS signaling to anticipate a DSA-REQ are beyond the scope of this standard.

In MR distributed scheduling case, RS shall check whether the QoS requirements can be supported on access link (RS ~ SS) or relay link (RS ~ RS) based on the received DSA-RSP. If it can be supported, RS shall send DSA-RSP to the MS or its subordinate RS and send the DSA-ACK to the BS. Otherwise, RS shall respond DSA-RSP with DSA-ACK with CC=reject to the MR-BS.

<u>The MR-BS checks whether the QoS requirements can be supported on relay link (MR-BS ~ RS). This process is illustrated in Table 125b.</u>

Table 125b – DSA initiated from SS through RS (Distributed scheduling case)

SS		<u>RS</u>		MR-BS
New service flow needed				
Check if resource are available				
Send DSA-REQ	DSA-REQ>	Receive / Send DSA-REQ	DSA-REQ>	Receive DSA-REQ
Set Timers T7 and T14				
Timer T14 Stops	<dsx-rvd< td=""><td>Receive / Send DSX-RVD</td><td><dsx-rvd< td=""><td>DSA-REQ integrity valid</td></dsx-rvd<></td></dsx-rvd<>	Receive / Send DSX-RVD	<dsx-rvd< td=""><td>DSA-REQ integrity valid</td></dsx-rvd<>	DSA-REQ integrity valid
				Check whether SS is authorized for Service
				Check whether service flow QoS can be supported on relay link (BS~RS)
				Create SFID
				If uplink AdmittedQoSParamSet is non-null, map service flow to CID
				If uplink ActiveQoSParamSet i non-null, Enable reception of data onew uplink service flow



^a Authorization happens prior to the DSA-REQ being received by the BS. The details of BS signaling to anticipate a DSA-REQ are beyond the scope of this standard.

6.3.14.9.3.2 BS-initiated DSA

Insert the following table the end of 6.3.14.9.3.2:

In MR centralized scheduling case, a RS only forwards each DSA messages from MR-BS to SS and vice versa. The MR-BS checks whether the QoS requirements can be supported both on relay link (MR-BS ~ RS) and on access link (RS-MS). This process is illustrated in Table 126a.

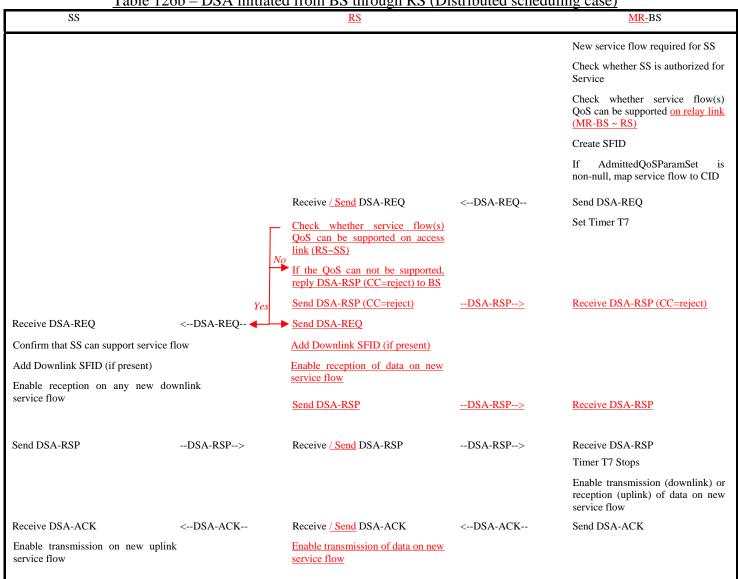
Table 126a – DSA initiated from BS through RS (Centralized scheduling case)

SS		RS		MR-BS
				New service flow required for SS
				Check whether SS is authorized for Service
				Check whether service flow(s) QoS can be supported both on relay link and on access link
				Create SFID
				If AdmittedQoSParamSet is non-null, map service flow to CID
Receive DSA-REQ	<dsa-req< td=""><td>Receive / Send DSA-REQ</td><td><dsa-req< td=""><td>Send DSA-REQ</td></dsa-req<></td></dsa-req<>	Receive / Send DSA-REQ	<dsa-req< td=""><td>Send DSA-REQ</td></dsa-req<>	Send DSA-REQ
				Set Timer T7
Confirm that SS can support serv	vice flow			
Add Downlink SFID (if present)				
Enable reception on any new service flow	downlink			
Send DSA-RSP	DSA-RSP>	Receive / Send DSA-RSP	DSA-RSP>	Receive DSA-RSP
				Timer T7 Stops
				Enable transmission (downlink) or reception (uplink) of data on new service flow
Receive DSA-ACK	<dsa-ack< td=""><td>Receive / Send DSA-ACK</td><td><dsa-ack< td=""><td>Send DSA-ACK</td></dsa-ack<></td></dsa-ack<>	Receive / Send DSA-ACK	<dsa-ack< td=""><td>Send DSA-ACK</td></dsa-ack<>	Send DSA-ACK
Enable transmission on new uservice flow	uplink			

In MR distributed scheduling case, the MR-BS only checks whether the QoS requirements can be supported on relay link (MR-BS ~ RS).

RS shall check whether the QoS requirements can be supported on access link (RS-SS) or relay link(RS-RS) based on the received DSA-REQ. If it can be supported, RS shall send DSA-REQ to the MS or its subordinate RS and send the DSA-RSP to the MR-BS. Otherwise, RS shall respond the DSA-REQ with DSA-RSP with CC=reject to the BS. This process is illustrated in Table 126b.

Table 126b – DSA initiated from BS through RS (Distributed scheduling case)



DSC 6.3.14.9.4

SS-initiated DSC 6.3.14.9.4.1

Insert the following table the end of 6.3.14.9.4.1:

In MR centralized scheduling case, a RS only forwards each DSC messages from SS to MR-BS and vice versa. The MR-BS checks whether the modified requirements can be supported both on relay link (MR-BS ~ RS) and on access link (RS-MS). This process is illustrated in Table 127a.

Table 127a – SS-initiated DSC through RS (Centralized scheduling case)

MR-BS		<u>RS</u>		SS
				Service flow requires modifying
Receive DSC-REQ	<dsc-req< th=""><th>Receive / Send DSC-REQ</th><th><dsc-req< th=""><th>Send DSC-REQ Set Timers T7 and T14</th></dsc-req<></th></dsc-req<>	Receive / Send DSC-REQ	<dsc-req< th=""><th>Send DSC-REQ Set Timers T7 and T14</th></dsc-req<>	Send DSC-REQ Set Timers T7 and T14
DSC-REQ integrity valid	DSX-RVD>	Receive / Send DSX-RVD	DSX-RVD>	Timer T14 Stops
Validate Request both on relay on access link	link and			
Modify service flow				
Increase Channel Bandwidth if	Required			
Send DSC-RSP	DSC-RSP>	Receive / Send DSC-RSP	DSC-RSP>	Receive DSC-RSP
				Timer T7 Stops
				Modify service flow
				Adjust Payload Bandwidth
Receive DSC-ACK	<dsc-ack< td=""><td>Receive / Send DSC-ACK</td><td><dsc-ack< td=""><td>Send DSC-ACK</td></dsc-ack<></td></dsc-ack<>	Receive / Send DSC-ACK	<dsc-ack< td=""><td>Send DSC-ACK</td></dsc-ack<>	Send DSC-ACK
Decease Channel Bandwidth if	Required			

In MR distributed scheduling case, RS shall check whether the modified requirements can be supported on access link (RS-SS) or relay link(RS-RS) based on the received DSA-RSP. If it can be supported, RS shall send DSA-RSP to the MS or its subordinate RS and send the DSA-ACK to the BS. Otherwise, RS shall respond DSA-RSP with DSA-ACK with CC=reject to the MR-BS.

The MR-BS checks whether the modified requirements can be supported on relay link (MR-BS ~ RS). This process is illustrated in Table 127b.

Table 127b – SS-initiated DSC through RS (Distributed scheduling case)

MR-BS		RS		SS
				Service flow requires modifying
Receive DSC-REQ	<dsc-req< td=""><td>Receive / Send DSC-REQ</td><td><dsc-req< td=""><td>Send DSC-REQ Set Timers T7 and T14</td></dsc-req<></td></dsc-req<>	Receive / Send DSC-REQ	<dsc-req< td=""><td>Send DSC-REQ Set Timers T7 and T14</td></dsc-req<>	Send DSC-REQ Set Timers T7 and T14
OSC-REQ integrity valid	DSX-RVD>	Receive / Send DSX-RVD	DSX-RVD>	Timer T14 Stops
Validate Request <u>on relay</u> (<u>MR-BS ~ SS</u>)	<u>link</u>			
Modify service flow				
Increase Channel Bandwidth if Rec	quired			
Send DSC-RSP	DSC-RSP>	Receive / Send DSC-RSP	DSC-RSP>	Receive DSC-RSP
		Check whether required service flow(s) QoS can be supported on access link (RS~SS)		
		If the QoS can not be supported, reply DSC-ACK (CC=reject) to SS	<u>No</u>	
Receive DSC-ACK (CC=reject)	<dsc-ack< td=""><td>Send DSC-ACK (CC=reject)</td><td><u>Yes</u></td><td></td></dsc-ack<>	Send DSC-ACK (CC=reject)	<u>Yes</u>	
		Send DSC-RSP	DSC-RSP>	Receive DSC-RSP
		Modify service flow		Timer T7 Stops
		Increase access uplink and		Modify service flow
		decrease relay downlink Payload Bandwidth if required.		Adjust Payload Bandwidth
	<dsc-ack< td=""><td>Receive / Send DSC-ACK</td><td><dsc-ack< td=""><td>Send DSC-ACK</td></dsc-ack<></td></dsc-ack<>	Receive / Send DSC-ACK	<dsc-ack< td=""><td>Send DSC-ACK</td></dsc-ack<>	Send DSC-ACK

Send DSC-ACK	DSC-ACK>	Receive DSC-ACK
		Modify service flow_
		Decrease access uplink and increase relay downlink Payload Bandwidth if required.

6.3.14.9.4.2 BS-initiated DSC

Insert the following table the end of 6.3.14.9.4.2:

In MR centralized scheduling case, a RS only forwards each DSC messages from MR-BS to SS and vice versa. The MR-BS checks whether the modified requirements can be supported both on relay link (MR-BS ~ RS) and on access link (RS ~ MS). This process is illustrated in Table 128a.

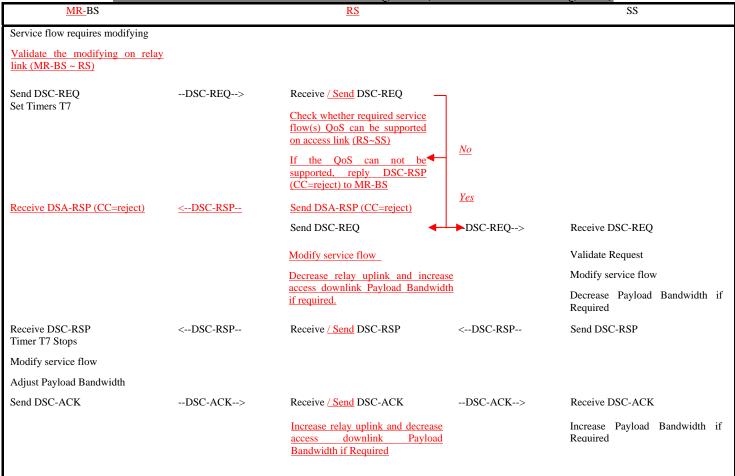
Table 128a – MR-BS-initiated DSC through RS (Centralized scheduling case)

MR-BS		<u>RS</u>		SS
Service flow requires modifying				
Validate the modifying both on relay link and on access link				
Send DSC-REQ	DSC-REQ>	Receive / Send DSC-REQ	DSC-REQ>	Receive DSC-REQ
Set Timers T7				Validate Request
				Modify service flow
				Decrease Payload Bandwidth if Required
Receive DSC-RSP Timer T7 Stops	<dsc-rsp< td=""><td>Receive / Send DSC-RSP</td><td><dsc-rsp< td=""><td>Send DSC-RSP</td></dsc-rsp<></td></dsc-rsp<>	Receive / Send DSC-RSP	<dsc-rsp< td=""><td>Send DSC-RSP</td></dsc-rsp<>	Send DSC-RSP
Modify service flow				
Adjust Payload Bandwidth				
Send DSC-ACK	DSC-ACK>	Receive / Send DSC-ACK	DSC-ACK>	Receive DSC-ACK
				Increase Payload Bandwidth if Required

<u>In MR distributed scheduling case, the MR-BS only checks whether the modified requirements can be supported on relay link (MR-BS ~ RS).</u>

And then the RS shall check whether the modified requirements can be supported on access link (RS-SS) or relay link (RS-RS) based on the received DSC-REQ. If the requirement can be supported, the RS transfers a DSC-REQ message to the MS or its subordinate RS and replies the DSC-RSP to the MR-BS. Otherwise, RS shall reply a DSC-RSP with CC=reject in order to inform the requirements can be not supported. This process is illustrated in Table 128b.

<u>Table 128b – MR-BS-initiated DSC through RS (Distributed scheduling case)</u>



6.3.14.9.5 Connection release 6.3.14.9.5.1 SS-initiated DSD

Insert the following table the end of 6.3.14.9.5.1:

<u>In MR centralized scheduling case, the MR-BS shall delete the service flow both on relay link (MR-BS ~ RS)</u> and on access link (RS-SS). This process is illustrated in Table 129b.

Table 129a – DSD-initiated from SS through RS (Centralized scheduling case)

SS		<u>RS</u>		MR-BS
Service flow no longer needed				
Delete service flow				
Send DSD-REQ	DSD-REQ>	Receive / Send DSD-REQ	DSD-REQ>	Receive DSD-REQ
				Verify SS is service flow "owner"
				Delete service flow both on relay link and on access link
Receive DSD-RSP	<dsd-rsp< th=""><th>Receive / Send DSD-RSP</th><th><dsd-rsp< th=""><th>Send DSD-RSP</th></dsd-rsp<></th></dsd-rsp<>	Receive / Send DSD-RSP	<dsd-rsp< th=""><th>Send DSD-RSP</th></dsd-rsp<>	Send DSD-RSP

In MR distributed scheduling case, the MR-BS shall delete the service flow on relay link (MR-BS ~ RS). And then the RS shall delete the service flow on access link (RS-SS). This process is illustrated in Table 128b.

Table 129b – DSD-initiated from SS through RS (Distributed scheduling case)

SS		<u>RS</u>		MR-BS
Service flow no longer needed				
Delete service flow				
Send DSD-REQ	DSD-REQ>	Receive / Send DSD-REQ	DSD-REQ>	Receive DSD-REQ
				Verify SS is service flow "owner"
				Delete service flow <u>on relay link</u> (MR-BS ~ RS)
Receive DSD-RSP	<dsd-rsp< td=""><td>Receive / Send DSD-RSP</td><td><dsd-rsp< td=""><td>Send DSD-RSP</td></dsd-rsp<></td></dsd-rsp<>	Receive / Send DSD-RSP	<dsd-rsp< td=""><td>Send DSD-RSP</td></dsd-rsp<>	Send DSD-RSP
		Delete service flow on access link (RS ~ SS)		

6.3.14.9.5.2 BS-initiated DSD

Insert the following table the end of 6.3.14.9.5.2:

<u>In MR centralized scheduling case, the MR-BS shall delete the service flow both on relay link (MR-BS ~ RS) and on access link (RS-SS).</u> This process is illustrated in Table 130b.

<u>Table 130a – DSD-initiated from MR-BS through RS (Centralized scheduling case)</u>

SS	1300 DSD initiated	RS RS	b (centralized sene	MR-BS
				Service flow no longer needed
				Delete service flow both on relay link and on access link
				Determine associated SS for this service flow
Receive DSD-REQ Delete service flow	<dsd-req< td=""><td>Receive / Send DSD-REQ</td><td><dsd-req< td=""><td>Send DSD-REQ</td></dsd-req<></td></dsd-req<>	Receive / Send DSD-REQ	<dsd-req< td=""><td>Send DSD-REQ</td></dsd-req<>	Send DSD-REQ
Send DSD-RSP	DSD-RSP>	Receive / Send DSD-RSP	DSD-RSP>	Receive DSD-RSP

In MR distributed scheduling case, the MR-BS shall delete the service flow on relay link (MR-BS ~ RS). And then the RS shall delete the service flow on access link (RS-SS). This process is illustrated in Table 130b.

<u>Table 130b – DSD-initiated from MR-BS through RS (Distributed scheduling case)</u>

SS		RS		MR-BS
				Service flow no longer needed
				Delete service flow on relay link (MR-BS ~ RS)
				Determine associated SS for this service flow
Receive DSD-REQ	<dsd-req< td=""><td>Receive / Send DSD-REQ</td><td><dsd-req< td=""><td>Send DSD-REQ</td></dsd-req<></td></dsd-req<>	Receive / Send DSD-REQ	<dsd-req< td=""><td>Send DSD-REQ</td></dsd-req<>	Send DSD-REQ
Delete service flow		Delete service flow on access lin (RS ~ SS)	<u>ık</u>	
Send DSD-RSP	DSD-RSP>	Receive / Send DSD-RSP	DSD-RSP>	Receive DSD-RSP

References

- [1] IEEE 802.16j-06/282, "Service flow management for RS", Kenji Saito, Takashi Inoue
- [2] IEEE 802.16j-07/188, "Shared Management Message: Format, Transfer and Security", Shulan Feng, Yanling Lu, Ting Li, Liangliang Zhang, Hisilicon Technologies.