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Re:	802.16REVe/D5 Sponsor Ballot	
Abstract	This contribution introduces a new compact-map type to provide a virtual multi-frame structure for periodic resource allocation.	
Purpose	Discuss and adopt the suggestion into P802.16e/D6.	
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A compact MAP message to provide a virtual multi-frame structure for a periodic fixed bandwidth assignment scheme

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1. Document Goal

This document introduces a new type of compact MAP message to facilitate a periodic resource allocation for the UGS-type of service classes, e.g., VoIP. By defining a virtual multi-frame structure, it shall allow for designating the fixed bandwidth assignment frames in an efficient manner. The objective of the current proposal is to reduce the MAP overhead inherent to the current specification, especially when the number of connections increases.

2. Problem

In the current specification, MAP message is required for dynamic bandwidth allocation, i.e., MAP message required for every frame in which the bandwidth is allocated. In general, it may suffer from an enormous overhead associated with the dynamic bandwidth allocation, especially when the number of connections increases. In particular, when a payload size is relatively small, the overhead problem becomes critical. One particular example is a VoIP service, which generates a small payload in a periodic manner, e.g., 20 byte payload generated every 20ms for G.729 codec. In fact, any type of CBR service is faced with the same problem. As all these types of services require a periodic resource allocation, it may not be necessary to resort to the MAP message in every frame. In other words, it is sufficient to notify an MSS of the bandwidth allocation information only once with the corresponding allocation period in terms of the number of frames. The same resource allocation information is used implicitly for the later frames until the session is over. Figure 1 illustrates the dynamic bandwidth assignment and fixed bandwidth assignment when a bandwidth reservation is required for every 4 frames. The obvious advantage of the fixed bandwidth assignment is that only a single MAP message is needed at the beginning of each session, which eliminates an overhead associated with individual MAP every time a bandwidth is reserved. In case that a channel condition is changed to incur the different PHY mode during a session, however, old MAP information is not valid any more, which makes the fixed assignment useless. To remedy this problem, new allocation information for the MSS subject to a PHY mode change can be included in the subsequent MAP message. Therefore, all MSS's subject to the fixed assignment still have to listen to all MAP messages in every frame, which makes sure that a new allocation is addressed. We note that the fixed bandwidth allocation can be implemented without changing the current specification, simply because the compact MAP identifies every CID allocated in the corresponding frame. As long as no allocation information for the MSS subject to the fixed assignment is

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found in the MAP message in a frame, it simply assumes that the previous allocation is still valid. In other words, the fixed assignment is maintained until a channel condition changes.

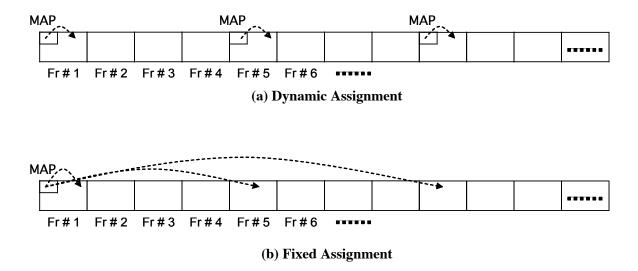


Fig 1. Dynamic Bandwidth Reservation Schemes for Periodic Allocation

In spite of the advantage of fixed assignment, its implementation is not straightforward in practice. The current specification addresses the bandwidth assignment in terms of the number of subchannels in the order of CID that appears in the MAP. In such a format, a region specified for the fixed assignment can be overwritten by another MSS, which makes the fixed assignment useless.

3. The Proposed Approach

To remedy the overwritten problem discussed in the previous section, we propose to use the *periodic* fixed bandwidth assignment scheme, which allows for refreshing the MAP message in a periodic manner. For a given period, a special update MAP message is used to announce that a currently designated region reserved for the fixed assignment must not be overwritten by all other MSS's while updating a change in the PHY mode if necessary. Whenever an MSS listens to the update MAP message transmitted in a given period, it finds the fixed assignment region, which must not be used by itself during that period.

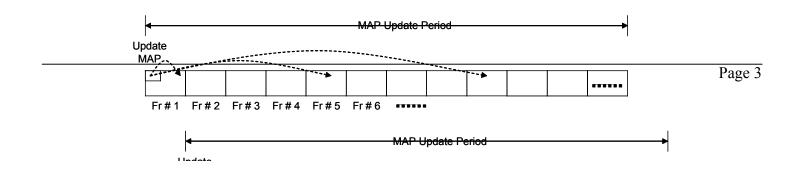
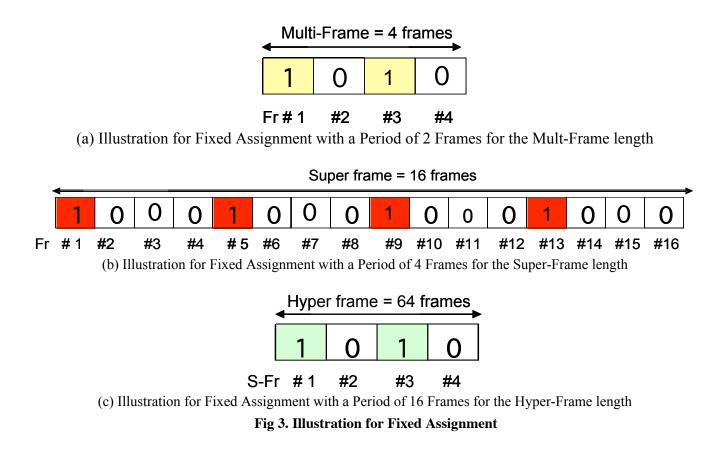


Fig 2. Periodic Fixed Bandwidth Assignment

Given a MAP update period for the fixed assignment, the update MAP message at the beginning of each period must indicate a specific frame in which a fixed assignment is applied per connection (See Figure 2). One possible way of indicating the frames of the fixed assignment is to use a bit map. In this case, as the update period increases, a size of bit map becomes excessive. In this contribution, we consider a notion of a multi-frame structure, which can be implemented in a virtual sense. As shown in Figure 3, three different types of frames are defined in a hierarchical structure: hyper-frame, super-frame, and multi-frame. Each multi-frame is composed of 4 frames. Each super-frame is composed of 4 multi-frames. Finally, each hyper-frame is composed of 4 super-frames. First of all, the update period is specified in terms of the different frame types. More specifically, the update period is given by a multiple of 4 frames, i.e., 4 frames (1 multi-frame), 8 frames (2 multi-frames), 16 frames (1 super-frame), 32 frames (2 super-frames), 64 frames (1 hyper-frame), and 128 frames (2 hyper-frames), which can be represented by a 3-bit long update period field in a MAP message. Once the update period is given, the fixed assignment is specified by the 3-bit long fixed assignment period filed and 3-bit long frame start offset field in a MAP message.

To accommodate the virtual frame structure addressed in the above, a new type of compact DL-MAP message is introduced, i.e., DL_MAP Type = 6 in Compact_DL-MAP_IE.

The difference from the existing compact DL-MAP is to include the update period and the fixed assignment period and start frame offset of assignment. Note that the fixed allocation period can take an integer in $[1,2^n]$, n = 1,2,3,...,7 under the current proposal. Figure 3 illustrates the proposed structure for the fixed assignment with various periods.



4. Specific Text Changes in the Standard

Remedy 1:

[Add the following tables to "the baseline document Section 6.3.2.3.43.6 Compact DL-MAP IE in page 118."] 6.3.2.3.43.6.7 Compact DL-MAP IE format for normal subchannel with multi-frame structure The format of Comapct DL-MAP IE for normal subchannel with multi-frame structure is presented in Table xx.

Syntax	Size	Notes
Compact_DL-MAP_IE () {		
DL-MAP Type =	3 bits	
UL-MAP append	1 bits	
RCID_IE	Variable	See Reduced CID section
Update Indicator	1 bits	1: Not updated for fixed assignment 0: Updated for fixed assignment
If (Update Indicator == 0) {		

		000: 1 Multi frame 001: 2 Multi frames
		010: 1 Super frame
		011: 2 Super frames
		100: 1 Hyper frame
Update type	3 bits	101: 2 Hyper frames
		110: 4 Hyper frames
		111: De-allocation
		The allocation ends at the update-type
		frames.
		The subchannels are allocated by the
Fixed Assignment Period (=p)	3 bits	period of every 2 [^] p frames.
		The terminal starts at the frame of
		which the number has the same 3 LSB
Frame Offset	3 bits	as the specified frame offset. If the
		current frame is specified, the terminal
		should start in 8 frames
Nep code	4 bits	Code of encoder packet bits
	4 0113	(see 8.4.2.3.5)
Nsch code	4 bits	Code of allocated subchannels
		(see 8.4.2.3.5)
Nband	Nb-Band bits	Number of bands, $0 =$ use BITMAP instead
if (Nband == 0) {		instead
Band BITMAP	Nb-BITMAP bits	n-th LSB is 1 if n-th band is selected
}else {		
for (i=0;i< Nband ; i++)		
Band Index	Nb-Index bits	Band selection.
}		Dwild Delevitori
,		Indicates the subchannel alocation
		mode.
		00 = same number of subchannels
		for the selected bands
Allocation Mode	2 bits	01 = different number of ubchannels
		for the selected bands
		10 = total number of subchannels
		for the selected bands determined
		by Nsch code and Nep code
Deserved	2 hita	10, 11 = reserved
$\frac{Reserved}{If (Allocation Mode = -00)}$	2 bits	Shall be set to zero
If (Allocation Mode == 00){ No. Subchannels	Q hita	
	8 bits	
} else if (Allocation Mode == 01) {		If Million die Orlander und im die mei 1
for (i-Oriz hand accent in 1)		If Nband is 0, band count is the number of '1' in Band BITMAP. Otherwise
for (i=0;i< band count ;i++) {		band count is Nband.
No. Subchannels	8 bits	
	0.0103	
}		
H-ARQ_Control_IE for UL	variable	

CQICH_Control_IE for UL	variable	
}		
if (UL-MAP append){		
Repetition Indicator	1 bits	1: Not updated for fixed assignment 0: Updated for fixed assignment
If (Update Indicator == 0) {		•
Update type	3 bits	000: 1 Multi frame 001: 2 Multi frames 010: 1 Super frame 011: 2 Super frames 100: 1 Hyper frame 101: 2 Hyper frames 110: 4 Hyper frames 111: Reserved
Fixed Assignment Period (=p)	3 bits	The subchannels are allocated by the
Frame Offset	3 bits	period of every 2^p frames. The terminal starts at the frame of which the number has the same 3 LSB as the specified frame offset. If the current frame is specified, the terminal should start in 8 frames
Nep code for UL	4 bits	Code of encoder packet bits (see 8.4.2.3.5)
Nsch code for UL	4 bits	Code of allocated subchannels (see 8.4.2.3.5)
Nband	Nb-Band bits	Number of bands, 0 = use BITMAP instead
if (Nband == 0) {		
Band BITMAP	Nb-BITMAP bits	n-th LSB is 1 if n-th band is selected
} else {		
for (i=0;i< Nband ; i++)		
Band Index	Nb-Index bits	Band selection.
}		
Allocation Mode	2 bits	Indicates the subchannel alocation mode. 00 = same number of subchannels for the selected bands 01 = different number of ubchannels for the selected bands 10 = total number of subchannels for the selected bands determined by Nsch code and Nep code 10, 11 = reserved
Reserved	2 bits	Shall be set to zero
If (Allocation Mode == 00){		
No. Subchannels	8 bits	
} else if (Allocation Mode == 01) {		
for (i=0;i< band count ;i++) {		If Nband is 0, band count is the number of '1' in Band BITMAP. Otherwise band count is Nband.

No. Subchannels	8 bits	
}		
}		
H-ARQ_Control_IE for UL	variable	
CQICH_Control_IE for UL	variable	
}		
}		

5. Simulation Results

Table 1 illustrates the MAP size reduced by the proposed approach for band AMC MAP IE when bandwidth is periodically allocated every 4 frames for N MSS's. While the MAP IE is generated at a rate of N MAP IE's/4 frames with the dynamic assignment, we assume that the MAP IE is updated at a period of 128 frames, i.e., MAP IE generation rate of N MAP IE's/128 frames. Assuming that No. band = 1 and N = 70, the existing MAP IE incurs an overhead of 16% for a frame structure with DL:UL = 24:12, while it can be reduced to about 4% with the proposed MAP IE. In general, the overhead reduction increases with the number of MSS's.

Table 1. Comparison of MAP size for the existing and proposed MAP IE: Band AMC MAP IE

	Existing MAP IE	Proposed MAP IE
MAP size: UL+ DL (A)	136 + 16 * No. band	148 + 16 * No. band
MAP IE Rate (B)	N MAP IE's/4 frames	N MAP IE's/128 frame
MAP Size (bits)/frame:	(136 + 16 * No. band) *	(148 + 16 * No. band) * N
(A) * (B)	N bits / 4 frames	bits / 128 frames

For the simulation of the proposed scheme, we made the system-level simulator based on IEEE 802.16d/d5. It is the TDD-OFDMA system with 10MHz bandwidth, total 1024 subcarriers, and a frame structure with DL:UL = 24:12.

We consider a multi-cell structure of 19 cells with 3 sectors, each with 1 km radius. We assume that all MSS's are uniformly distributed in each cell. Data rate for each MSS is determined by the AMC scheme, following the SINR requirement identified by a link-level simulation. Using the COST-231 model for a path loss, i.e., for BS

height of 32m and MS height of 1.5m, the loss at distance d [km] is given as follows [4]:

Log-normal shadowing is considered for a large-scale fading model, i.e., shadowing is modeled by , where

 Z_k is a Gaussian random variable ~ N(0,1) and $\delta = 8$ dB. Furthermore, we assume that inter-cell interference is fully loaded, i.e., all subcarriers of each subcarriers are all used. Furthermore, we assume the error-free transmission, i.e., no ARQ protocol invoked.

We consider three different types of traffic sources, i.e., voice, ethernet, and video sources. Each of these traffic sources are characterized by the parameters shown in Table 2. For a video source, we use the video streaming traffic in [3]. The voice and video traffic is subject to delay constraints. For a voice source, both downlink and uplink is considered. In the current analysis, the highest priority is given to the voice traffic, i.e., bandwidth scheduled for them ahead of the other traffic classes. The remaining resource is allocated for video real-time traffic. Finally, the ethernet traffic is served only with the remaining resource. To decide which connection to serve among the same traffic class, the proportional fairness (PF) scheduling algorithm is applied.

The proposed fixed allocation scheme is applied just at voice and video traffic and is not applied at ethernet traffic. We consider the just diversity channel MAP in the result of MAP symbol overhead ratio, because using the only diversity subchannel.

	Voice	Video	Ethernet
Arrival rate	Constant	10 fps	Pareto ($\lambda = 75$ packtes/sec)
Packet length	16 kbps	Data rate : 32 kbps	$\mu = 2800$ bits/packet
Max delay	0.04 sec	-	-

We simulated when fixed allocation scheme is not used and when fixed allocation scheme is used by the update period with 128, 64, and 16 frames. In Figure 4, if the fixed allocation scheme is used, the map overhead symbol ratio of those schemes is much lower than that of general MAP allocation scheme. The MAP overhead symbol ratio of general MAP allocation scheme is about $14 \sim 18\%$. The MAP overhead symbol ratio of the proposed fixed allocation scheme is about $1 \sim 4\%$. Therefore, the gap of MAP symbol ratio between two schemes is about $13 \sim 16\%$.

In Figure 5, if the fixed allocation scheme is used, the system throughput of those schemes is much higher than that of general MAP allocation scheme. The maximal throughput using the proposed fixed allocation scheme is about 2.7Mbps. But, the maximal throughput using the general allocation scheme is about 2Mbps. Therefore, the gap of throughput between two schemes is about 30%. After all, decreasing the total amount of MAPs, we will increase the system throughput. According to the update period, the throughput and MAP overhead symbol ratio are similar. Because this simulation uses the only diversity channel MAP, the total amount of the MAP symbol using both 128 and 16 update period is similar and those throughput is similar, too. In Figure 6, we show the throughput according to the allocation period. The lager the allocation period is, the bigger the throughput is. Instead of increasing the throughput, the packet drop ratio of the real-time traffic will somewhat increase.

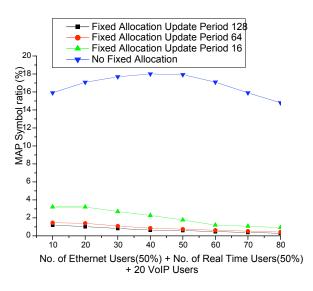


Fig 4. Average MAP symbol ratio

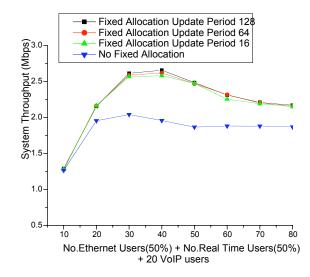


Fig 5. Average throughput

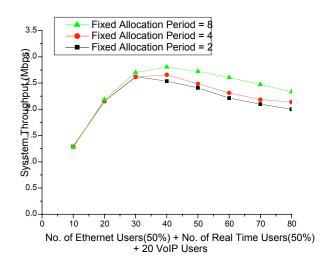


Fig 6. Average throughput by fixed allocation period

6. References

[1] IEEE P802.16e/D3-2004 Amendment for Physical and Medium Access Control Layers for Combined Fixed and Mobile Operation in Licensed Band.

[2] IEEE P802.16-REVd/D5-2004 Air Interface for Fixed Broadband Wireless Access Systems

[3] 3GPP-2, "1xEV-DV Evaluation Methodology Addendum (V6)," July 25, 2001.