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| Abstract | Defines the missing details on per stream power control feedback in CQICH enhanced allocation IE | |
| Purpose | Adoption of proposed changes into P802.16e | |
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Per Stream Power Control in CQICH Enhanced Allocation IE

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1 Introduction

In section 8.4.5.4.15 of IEEE 80216e/D6, a few new types of feedback information are specified in CQICH_Enhanced_Alloc_IE() so that the SS can transmit feedback information of that specified type on the assigned CQICH. However, there is no detail on how the payload bits should be if the feedback type is “101”, i.e., when per stream power control feedback is required by the base.

This contribution provides the missing specifications on how per-stream power weighting feedback is performed. Power weighting on different streams can be very effective in dealing with the different spatial channel quality in MIMO communications. When the power weighting of streams are ordered and summed up to one, the quantization of the power weightings can be done with very fine granularity and also very efficiently (three bits for up to two streams, six bits for up to three streams, and nine bits for up to four streams).

2 Feedback Types and Per Stream Power Control Feedback

In table 302a of section 8.4.5.4.15 of IEEE 80216e/D6, a three bit feedback type field is defined for each CQICH. For example, “000” indicates that the CQICH should carry the information about Fast DL measurement/Default Feedback with antenna grouping, i.e., both the DL SNR measurement and MIMO mode can be fed back where some codewords for MIMO mode feedback are interpreted differently in case of antenna grouping (“000”), antenna selection (“001”), or reduced precoding matrix code book (“010”), as detailed in section 8.4.5.4.10.7. Feedback type “011” indicates that the MSS shall report the quantized MIMO precoding coefficient according to the mapping defined in section 8.4.5.4.10.6. Feedback type “100” indicates that the MSS shall report the index to the precoding matrix defined in 8.4.5.4.11. However, **bits “101” is define twice:** once for “channel matrix information” and once for “per stream power control”. The “channel matrix information” might mean the same as quantized MIMO coefficients by bits “011”. If so, it can be deleted. Otherwise, a new bit word out of the reserved ones must be used for “per stream power control”. **The encoding of the payload bits for “per stream power control” is not defined anywhere.**

Per stream power weighting can be very effective in dealing with the different spatial channel quality in MIMO communications. One way to inform the base about the different quality of spatial channels is to use multiple CQICH to convey the post-processing SNR of each stream. Then the BS has to choose a different coding/modulation to take advantage of that information. Sometime, there is no coding/modulation combination that can match the SNR of the stream, in which case adjusting the power of each stream can be found useful. Still, per stream SNR feedback and coding/modulation selection means additional control overhead for both uplink and downlink. In addition, if the receiver uses some non-linear MIMO demodulation methods such as ML or successive cancellation, there is no effective SNR for each stream to be computed. One alternative to reduce the overhead is to use the same code and modulation for all streams, but with different power weighting so that the mean squared error at the particular receiver (e.g., one uses successive cancellation) is equalized on each stream. Reference [1] gives a description of how the power weighting may be calculated. There are other cases that power weighting can be found useful too, such as described before. In addition, the power weighting can be combined with beamforming codebook weights.

For the quantization of power weightings for all streams, a more efficient quantization method is described here after recognizing the fact that the range of each stream can be refined after the power weightings of previous streams are quantized. The method sequentially quantizes the power weightings of the data streams in a numerical range that depends on the power weighting of the previously quantized stream powers. Also noted here that the streams are indexed in the order of decreasing power weighting and all power weightings sum up to one. So the number of bits assigned to quantize each stream can be smaller due to the decreasing range. Due to the fact that only a minimum dynamic range is quantized, the quantization granularity is very fine. In most of the cases a granularity of less than 0.01 in power difference is achieved with extreme cases being about 0.08.

The quantization scheme is given as (note that the quantization step includes both the lower and upper range values and the remaining B-2 levels are uniformly positioned between the lower and upper limits):

1. Determine the maximum number of data streams that the BS and MSS support based on the number of antennas at MSS and BS (e.g., up to N_s data streams)
2. Determine the power weighting P_1 of the strongest data stream between $1/N_s$ and 1. Quantize the squared root of P_1 (i.e., $\alpha_1 = \sqrt{P_1}$) with B_1 bits ($B_1=3$ if $N_s=2$, $B_1=4$ if $N_s=3$ or 4).
3. For the m -th stream where $m=2$ to N_s-1 , determine the power weighting of the m -th data stream that should be in the range of $\frac{1}{N_s+1-m} (1 - \sum_{n=1}^{m-1} \alpha_n^2) \leq \alpha_m^2 \leq \min(\alpha_{m-1}^2, 1 - \sum_{n=1}^{m-1} \alpha_n^2)$. Quantize α_m with B_m bits ($B_2=2$ if $N_s=3$, $B_2=3$ and $B_3=2$ if $N_s=4$).
4. The power weighting of the last stream is $1 - \sum_{n=1}^{N_s-1} \alpha_n^2$, which does not need to be fed back.

Thus the total amount of feedback (in number of bits) needed for the power weight is $\sum_{m=1}^{N_s-1} B_m$. A three-bit

CQICH is allocated if the feedback for up to $N_s=2$ streams is requested by the BS. A six-bit CQICH is allocated (or two three-bit CQICH) if the feedback for up to $N_s=3$ streams is requested by the BS. A six-bit and a three-bit CQICH (or 3 three-bit CQICH) are allocated if the feedback for up to $N_s=4$ streams is requested by the BS. Note that if the MSS preferred a stream number smaller than N_s , the remaining streams will be allocated with zero power.

3 Specific Text Changes

[NOTE: In the following, we propose that the working group adopt one of two functionally equivalent remedies that enable the per-stream power control scheme described above: The first version (Section 3.1) specifies the per-stream power control strategy with both equations and tables to show the specific encoding values. The second version (Section 3.2) omits the tables and contains only the equations.]

3.1 Remedy 1 – Version with Formulas and Tables

[Insert the following after Section 8.4.5.4.10.10:]

8.4.5.4.10.11 Per Stream Power Control

When the feedback type field in CQICH Enhanced Allocation IE is “101” = Per stream power control, the BS require the power weighting of each spatial streams that can be supported by the MSS if the BS considers sending more than one stream to this MSS. If required by the BS, the MSS shall report the square root of the power weighting factors of the spatial streams (i.e., to report α_i with $i=1..N_s$ (number of streams) where $\sum_i \alpha_i^2=1$). The first stream shall correspond to the largest weighting and the second stream to the second largest weighting, and so on. The power weighting of the last stream can be derived as the remaining power and thus needs not to be reported.

The feedback allocation and power weighting quantization procedure is as follows:

If the BS wants the MSS to feed back the power weightings for up to $N_s=2$ streams, one 3-bit CQI channel is allocated. A numerical range of $[\sqrt{1/2}, 1]$ is first uniformly divided into $2^3=8$ levels (i.e., with the interval between levels being $(1-\sqrt{1/2})/7$) and the MSS quantizes the squared root of the power weighting of the first stream to the nearest level according to table XXX. All “1” means the first stream uses all transmit power (i.e., a single stream is preferred by the MSS, rather than two streams).

Table XXX. Encoding of the payload bits for per-stream power control (up to 2 stream case)

| α_1 bits | 000 | 001 | 010 | 011 | 100 | 101 | 110 | 111 |
|-----------------|--------|--------|--------|--------|--------|--------|--------|--------|
| α_1 | 0.7071 | 0.7489 | 0.7908 | 0.8326 | 0.8745 | 0.9163 | 0.9582 | 1.0000 |

If the BS wants the MSS to feed back power weighting for up to $N_s=3$ streams, one 6-bit CQI (or two 3-bit CQI channels) is allocated. The first power weighting is quantized using 4 bits and the second using 2 bits. A numerical range of $[\sqrt{1/3}, 1]$ is first uniformly divided into $2^4=16$ levels (i.e., with the interval between levels being $(1-\sqrt{1/3})/15$) and the MSS quantizes the squared root of the power weighting of the first stream to the nearest level (denoted as α_1). Then, 2 bits are used to quantize the range of $[\frac{1}{2}(1-\alpha_1^2), \min(\alpha_1^2, 1-\alpha_1^2)]$. The exact quantization is given in table YYY-1 and YYY-2.

Table YYY-1. Encoding of the payload bits for per-stream power control (up to 3 stream case)

| α_1 bits | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 |
|-----------------|--------|--------|--------|--------|--------|--------|--------|--------|
| α_1 | 0.5774 | 0.6055 | 0.6337 | 0.6619 | 0.6901 | 0.7182 | 0.7464 | 0.7746 |

| | | | | | | | | |
|-----------------|--------|--------|--------|--------|--------|--------|--------|--------|
| | | | | | | | | |
| α_1 bits | 1000 | 1001 | 1010 | 1011 | 1100 | 1101 | 1110 | 1111 |
| α_1 | 0.8028 | 0.8309 | 0.8591 | 0.8873 | 0.9155 | 0.9436 | 0.9718 | 1.0000 |

Table YYY-2. Encoding of the payload bits for per-stream power control (up to 3 stream case)

| | | | | | | | | |
|-------------------|--------|--------|--------|--------|--------|--------|--------|--------|
| α_1 bits | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 |
| α_2 : "00" | 0.5774 | 0.5627 | 0.5470 | 0.5301 | 0.5118 | 0.4920 | 0.4706 | 0.4472 |
| "01" | 0.5774 | 0.5770 | 0.5759 | 0.5740 | 0.5712 | 0.5599 | 0.5355 | 0.5090 |
| "10" | 0.5774 | 0.5913 | 0.6048 | 0.6179 | 0.6306 | 0.6279 | 0.6005 | 0.5707 |
| "11" | 0.5774 | 0.6055 | 0.6337 | 0.6619 | 0.6901 | 0.6958 | 0.6655 | 0.6325 |
| | | | | | | | | |
| α_1 bits | 1000 | 1001 | 1010 | 1011 | 1100 | 1101 | 1110 | 1111 |
| α_2 : "00" | 0.4216 | 0.3934 | 0.3619 | 0.3261 | 0.2845 | 0.2340 | 0.1667 | 0 |
| "01" | 0.4799 | 0.4477 | 0.4118 | 0.3711 | 0.3238 | 0.2663 | 0.1897 | 0 |
| "10" | 0.5381 | 0.5020 | 0.4618 | 0.4162 | 0.3631 | 0.2986 | 0.2127 | 0 |
| "11" | 0.5963 | 0.5564 | 0.5118 | 0.4612 | 0.4024 | 0.3310 | 0.2357 | 0 |

If the BS wants the MSS to feed back power weighting for up to $N_s=4$ streams, one 6-bit CQI and one 3-bit CQI (or three 3-bit CQI channels) are allocated. The first power weighting is quantized using 4 bits, the second using 3 bits, and the third using 2 bits. A numerical range of $[\sqrt{1/4}, 1]$ is first uniformly divided into $2^4=16$ levels (i.e., with the interval between levels being $(1-\sqrt{1/4})/15$) and the MSS quantizes the squared root of the power weighting of the first stream to the nearest level (denoted as α_1). Then, 3 bits are used to quantize the range of $[\frac{1}{3}(1-\alpha_1^2), \min(\alpha_1^2, 1-\alpha_1^2)]$ for the second stream squared-root power weighting (denotes as α_2). Finally, 2 bits are used to quantize the range of $[\frac{1}{2}(1-\alpha_1^2-\alpha_2^2), \min(\alpha_2^2, 1-\alpha_1^2-\alpha_2^2)]$ for the third stream squared-root power weighting. The exact quantization is given in table ZZZ-1, ZZZ-2, and ZZZ-3.

Table ZZZ-1. Encoding of the payload bits for per-stream power control (up to 4 streams case)

| | | | | | | | | |
|-----------------|--------|--------|--------|--------|--------|--------|--------|--------|
| α_1 bits | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 |
| α_1 | 0.5000 | 0.5333 | 0.5667 | 0.6000 | 0.6333 | 0.6667 | 0.7000 | 0.7333 |
| | | | | | | | | |
| α_1 bits | 1000 | 1001 | 1010 | 1011 | 1100 | 1101 | 1110 | 1111 |
| α_1 | 0.7667 | 0.8000 | 0.8333 | 0.8667 | 0.9000 | 0.9333 | 0.9667 | 1.0000 |

Table ZZZ-2. Encoding of the payload bits for per-stream power control (up to 4 streams case)

| | | | | | | | | |
|--------------------|--------|--------|--------|--------|--------|--------|--------|--------|
| α_1 bits | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 |
| α_2 : "000" | 0.5000 | 0.4884 | 0.4757 | 0.4619 | 0.4468 | 0.4303 | 0.4123 | 0.3925 |
| "001" | 0.5000 | 0.4948 | 0.4887 | 0.4816 | 0.4734 | 0.4641 | 0.4534 | 0.4336 |

| | | | | | | | | |
|--------------------|--------|--------|--------|--------|--------|--------|--------|--------|
| | | | | | | | | |
| “010” | 0.5000 | 0.5012 | 0.5017 | 0.5013 | 0.5001 | 0.4979 | 0.4945 | 0.4746 |
| “011” | 0.5000 | 0.5076 | 0.5147 | 0.5211 | 0.5267 | 0.5316 | 0.5356 | 0.5157 |
| “100” | 0.5000 | 0.5141 | 0.5277 | 0.5408 | 0.5534 | 0.5654 | 0.5767 | 0.5567 |
| “101” | 0.5000 | 0.5205 | 0.5407 | 0.5605 | 0.5800 | 0.5991 | 0.6178 | 0.5978 |
| “110” | 0.5000 | 0.5269 | 0.5537 | 0.5803 | 0.6067 | 0.6329 | 0.6589 | 0.6388 |
| “111” | 0.5000 | 0.5333 | 0.5667 | 0.6000 | 0.6333 | 0.6667 | 0.7000 | 0.6799 |
| | | | | | | | | |
| α_1 bits | 1000 | 1001 | 1010 | 1011 | 1100 | 1101 | 1110 | 1111 |
| α_2 : “000” | 0.3707 | 0.3464 | 0.3191 | 0.2880 | 0.2517 | 0.2073 | 0.1478 | 0 |
| “001” | 0.4095 | 0.3826 | 0.3525 | 0.3182 | 0.2780 | 0.2290 | 0.1633 | 0 |
| “010” | 0.4482 | 0.4189 | 0.3859 | 0.3483 | 0.3043 | 0.2506 | 0.1787 | 0 |
| “011” | 0.4870 | 0.4551 | 0.4193 | 0.3784 | 0.3306 | 0.2723 | 0.1942 | 0 |
| “100” | 0.5257 | 0.4913 | 0.4526 | 0.4085 | 0.3569 | 0.2940 | 0.2097 | 0 |
| “101” | 0.5645 | 0.5275 | 0.4860 | 0.4386 | 0.3833 | 0.3157 | 0.2251 | 0 |
| “110” | 0.6033 | 0.5638 | 0.5194 | 0.4688 | 0.4096 | 0.3373 | 0.2406 | 0 |
| “111” | 0.6420 | 0.6000 | 0.5528 | 0.4989 | 0.4359 | 0.3590 | 0.2560 | 0 |

Table ZZZ-3. Encoding of the payload bits for per-stream power control (up to 4 streams case)

| | | | | | | | | |
|-------------------------|---------|---------|---------|---------|---------|---------|---------|---------|
| $\alpha_1\alpha_2$ bits | 0000000 | 0000001 | 0000010 | 0000011 | 0000100 | 0000101 | 0000110 | 0000111 |
| α_3 : “00” | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 |
| “01” | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 |
| “10” | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 |
| “11” | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 | 0.5000 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 0001000 | 0001001 | 0001010 | 0001011 | 0001100 | 0001101 | 0001110 | 0001111 |
| α_3 : “00” | 0.4884 | 0.4851 | 0.4818 | 0.4785 | 0.4750 | 0.4715 | 0.4679 | 0.4643 |
| “01” | 0.4884 | 0.4884 | 0.4883 | 0.4882 | 0.4880 | 0.4878 | 0.4876 | 0.4873 |
| “10” | 0.4884 | 0.4916 | 0.4948 | 0.4979 | 0.5011 | 0.5042 | 0.5073 | 0.5103 |

| | | | | | | | | |
|-------------------------|---------|---------|---------|---------|---------|---------|---------|---------|
| | | | | | | | | |
| “11” | 0.4884 | 0.4948 | 0.5012 | 0.5076 | 0.5141 | 0.5205 | 0.5269 | 0.5333 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 0010000 | 0010001 | 0010010 | 0010011 | 0010100 | 0010101 | 0010110 | 0010111 |
| α_3 : “00” | 0.4757 | 0.4691 | 0.4622 | 0.4550 | 0.4475 | 0.4396 | 0.4315 | 0.4230 |
| “01” | 0.4757 | 0.4756 | 0.4753 | 0.4749 | 0.4742 | 0.4733 | 0.4722 | 0.4709 |
| “10” | 0.4757 | 0.4822 | 0.4885 | 0.4948 | 0.5009 | 0.5070 | 0.5129 | 0.5188 |
| “11” | 0.4757 | 0.4887 | 0.5017 | 0.5147 | 0.5277 | 0.5407 | 0.5537 | 0.5667 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 0011000 | 0011001 | 0011010 | 0011011 | 0011100 | 0011101 | 0011110 | 0011111 |
| α_3 : “00” | 0.4619 | 0.4517 | 0.4408 | 0.4292 | 0.4169 | 0.4036 | 0.3894 | 0.3742 |
| “01” | 0.4619 | 0.4617 | 0.4610 | 0.4598 | 0.4582 | 0.4559 | 0.4432 | 0.4258 |
| “10” | 0.4619 | 0.4716 | 0.4812 | 0.4905 | 0.4995 | 0.5082 | 0.4969 | 0.4775 |
| “11” | 0.4619 | 0.4816 | 0.5013 | 0.5211 | 0.5408 | 0.5605 | 0.5507 | 0.5292 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 0100000 | 0100001 | 0100010 | 0100011 | 0100100 | 0100101 | 0100110 | 0100111 |
| α_3 : “00” | 0.4468 | 0.4329 | 0.4176 | 0.4009 | 0.3825 | 0.3622 | 0.3397 | 0.3145 |
| “01” | 0.4468 | 0.4464 | 0.4451 | 0.4428 | 0.4353 | 0.4123 | 0.3866 | 0.3579 |
| “10” | 0.4468 | 0.4599 | 0.4726 | 0.4848 | 0.4882 | 0.4623 | 0.4335 | 0.4013 |
| “11” | 0.4468 | 0.4734 | 0.5001 | 0.5267 | 0.5410 | 0.5123 | 0.4804 | 0.4447 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 0101000 | 0101001 | 0101010 | 0101011 | 0101100 | 0101101 | 0101110 | 0101111 |
| α_3 : “00” | 0.4303 | 0.4124 | 0.3922 | 0.3694 | 0.3434 | 0.3135 | 0.2784 | 0.2357 |
| “01” | 0.4303 | 0.4296 | 0.4274 | 0.4204 | 0.3909 | 0.3568 | 0.3168 | 0.2682 |
| “10” | 0.4303 | 0.4469 | 0.4626 | 0.4714 | 0.4383 | 0.4001 | 0.3552 | 0.3008 |
| “11” | 0.4303 | 0.4641 | 0.4979 | 0.5224 | 0.4857 | 0.4434 | 0.3937 | 0.3333 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 0110000 | 0110001 | 0110010 | 0110011 | 0110100 | 0110101 | 0110110 | 0110111 |
| α_3 : “00” | 0.4123 | 0.3901 | 0.3643 | 0.3340 | 0.2978 | 0.2533 | 0.1947 | 0.1000 |
| “01” | 0.4123 | 0.4112 | 0.4077 | 0.3801 | 0.3390 | 0.2883 | 0.2216 | 0.1138 |
| “10” | 0.4123 | 0.4323 | 0.4511 | 0.4262 | 0.3801 | 0.3232 | 0.2485 | 0.1276 |
| “11” | 0.4123 | 0.4534 | 0.4945 | 0.4724 | 0.4212 | 0.3582 | 0.2754 | 0.1414 |

| | | | | | | | | |
|-------------------------|---------|---------|---------|---------|---------|---------|---------|---------|
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 0111000 | 0111001 | 0111010 | 0111011 | 0111100 | 0111101 | 0111110 | 0111111 |
| α_3 : "00" | 0.3925 | 0.3703 | 0.3442 | 0.3133 | 0.2759 | 0.2290 | 0.1645 | 0 |
| "01" | 0.3925 | 0.3914 | 0.3877 | 0.3566 | 0.3140 | 0.2606 | 0.1872 | 0 |
| "10" | 0.3925 | 0.4125 | 0.4312 | 0.3998 | 0.3521 | 0.2923 | 0.2099 | 0 |
| "11" | 0.3925 | 0.4336 | 0.4746 | 0.4431 | 0.3902 | 0.3239 | 0.2327 | 0 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1000000 | 1000001 | 1000010 | 1000011 | 1000100 | 1000101 | 1000110 | 1000111 |
| α_3 : "00" | 0.3707 | 0.3497 | 0.3251 | 0.2959 | 0.2606 | 0.2163 | 0.1554 | 0 |
| "01" | 0.3707 | 0.3696 | 0.3661 | 0.3367 | 0.2966 | 0.2461 | 0.1768 | 0 |
| "10" | 0.3707 | 0.3895 | 0.4072 | 0.3776 | 0.3325 | 0.2760 | 0.1983 | 0 |
| "11" | 0.3707 | 0.4095 | 0.4482 | 0.4184 | 0.3685 | 0.3059 | 0.2197 | 0 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1001000 | 1001001 | 1001010 | 1001011 | 1001100 | 1001101 | 1001110 | 1001111 |
| α_3 : "00" | 0.3464 | 0.3268 | 0.3038 | 0.2765 | 0.2435 | 0.2021 | 0.1452 | 0 |
| "01" | 0.3464 | 0.3454 | 0.3421 | 0.3147 | 0.2771 | 0.2300 | 0.1652 | 0 |
| "10" | 0.3464 | 0.3640 | 0.3805 | 0.3528 | 0.3108 | 0.2579 | 0.1853 | 0 |
| "11" | 0.3464 | 0.3826 | 0.4189 | 0.3910 | 0.3444 | 0.2858 | 0.2053 | 0 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1010000 | 1010001 | 1010010 | 1010011 | 1010100 | 1010101 | 1010110 | 1010111 |
| α_3 : "00" | 0.3191 | 0.3011 | 0.2799 | 0.2547 | 0.2244 | 0.1862 | 0.1338 | 0 |
| "01" | 0.3191 | 0.3182 | 0.3152 | 0.2899 | 0.2553 | 0.2119 | 0.1522 | 0 |
| "10" | 0.3191 | 0.3354 | 0.3505 | 0.3251 | 0.2863 | 0.2376 | 0.1707 | 0 |
| "11" | 0.3191 | 0.3525 | 0.3859 | 0.3602 | 0.3173 | 0.2633 | 0.1892 | 0 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1011000 | 1011001 | 1011010 | 1011011 | 1011100 | 1011101 | 1011110 | 1011111 |
| α_3 : "00" | 0.2880 | 0.2717 | 0.2526 | 0.2299 | 0.2025 | 0.1680 | 0.1207 | 0 |
| "01" | 0.2880 | 0.2872 | 0.2845 | 0.2616 | 0.2304 | 0.1913 | 0.1374 | 0 |
| "10" | 0.2880 | 0.3027 | 0.3164 | 0.2934 | 0.2584 | 0.2145 | 0.1541 | 0 |
| "11" | 0.2880 | 0.3182 | 0.3483 | 0.3251 | 0.2864 | 0.2377 | 0.1707 | 0 |

| | | | | | | | | |
|-------------------------|---------|---------|---------|---------|---------|---------|---------|---------|
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1100000 | 1100001 | 1100010 | 1100011 | 1100100 | 1100101 | 1100110 | 1100111 |
| α_3 : "00" | 0.2517 | 0.2374 | 0.2207 | 0.2009 | 0.1769 | 0.1468 | 0.1055 | 0 |
| "01" | 0.2517 | 0.2509 | 0.2486 | 0.2286 | 0.2013 | 0.1671 | 0.1200 | 0 |
| "10" | 0.2517 | 0.2645 | 0.2764 | 0.2563 | 0.2258 | 0.1874 | 0.1346 | 0 |
| "11" | 0.2517 | 0.2780 | 0.3043 | 0.2841 | 0.2502 | 0.2076 | 0.1492 | 0 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1101000 | 1101001 | 1101010 | 1101011 | 1101100 | 1101101 | 1101110 | 1101111 |
| α_3 : "00" | 0.2073 | 0.1955 | 0.1818 | 0.1654 | 0.1457 | 0.1209 | 0.0869 | 0 |
| "01" | 0.2073 | 0.2067 | 0.2047 | 0.1883 | 0.1658 | 0.1376 | 0.0989 | 0 |
| "10" | 0.2073 | 0.2178 | 0.2277 | 0.2111 | 0.1859 | 0.1543 | 0.1109 | 0 |
| "11" | 0.2073 | 0.2290 | 0.2506 | 0.2340 | 0.2061 | 0.1710 | 0.1229 | 0 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1110000 | 1110001 | 1110010 | 1110011 | 1110100 | 1110101 | 1110110 | 1110111 |
| α_3 : "00" | 0.1478 | 0.1395 | 0.1296 | 0.1180 | 0.1039 | 0.0862 | 0.0620 | 0 |
| "01" | 0.1478 | 0.1474 | 0.1460 | 0.1343 | 0.1183 | 0.0982 | 0.0705 | 0 |
| "10" | 0.1478 | 0.1553 | 0.1624 | 0.1506 | 0.1326 | 0.1101 | 0.0791 | 0 |
| "11" | 0.1478 | 0.1633 | 0.1787 | 0.1669 | 0.1470 | 0.1220 | 0.0876 | 0 |
| | | | | | | | | |
| $\alpha_1\alpha_2$ bits | 1111000 | 1111001 | 1111010 | 1111011 | 1111100 | 1111101 | 1111110 | 1111111 |
| α_3 : "00" | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| "01" | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| "10" | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| "11" | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

[Add a new section 11.7.8.11]

11.7.8.11 Advanced Receiver Capability

This field indicates whether the MSS is advanced receiver capable

| Type | Length | Value | Scope |
|------|--------|-------|-------|
|------|--------|-------|-------|

| | | | |
|----|---|---|--------------------|
| 21 | 1 | Bit 0: Successive Interference Receiver Capability Bit 1-7: Reserved | REG-REQ REG-RSP |
|----|---|---|--------------------|

----- End of Text Changes for Remedy 1 -----

3.2 Remedy 2 – Version with formulas only (no tables)

[Insert the following after Section 8.4.5.4.10.10:]

8.4.5.4.10.11 Per Stream Power Control

When the feedback type field in CQICH Enhanced Allocation IE is “101” = Per stream power control, the BS require the power weighting of each spatial streams that can be supported by the MSS if the BS considers sending more than one stream to this MSS. If required by the BS, the MSS shall report the square root of the power weighting factors of the spatial streams (i.e., to report α_i with $i=1..N_s$ (number of streams) where $\sum_i \alpha_i^2=1$). The first stream shall correspond to the largest weighting and the second stream to the second largest weighting, and so on. The power weighting of the last stream can be derived as the remaining power and thus needs not to be reported.

The feedback allocation and power weighting quantization procedure is as follows:

If the BS wants the MSS to feed back the power weightings for up to $N_s=2$ streams, one 3-bit CQI channel is allocated. A numerical range of $[\sqrt{1/2}, 1]$ is first uniformly divided into $2^3=8$ levels (i.e., with the interval between levels being $(1-\sqrt{1/2})/7$) and the MSS quantizes the squared root of the power weighting of the first stream to the nearest level. The value of “1” means the first stream uses all transmit power (i.e., a single stream is preferred by the MSS, rather than two streams).

If the BS wants the MSS to feed back power weighting for up to $N_s=3$ streams, one 6-bit CQI (or two 3-bit CQI channels) is allocated. The first power weighting is quantized using 4 bits and the second using 2 bits. A numerical range of $[\sqrt{1/3}, 1]$ is first uniformly divided into $2^4=16$ levels (i.e., with the interval between levels being $(1-\sqrt{1/3})/15$) and the MSS quantizes the squared root of the power weighting of the first stream to the nearest level (denoted as α_1). Then, 2 bits are used to quantize the range of $[\frac{1}{2}(1-\alpha_1^2), \min(\alpha_1^2, 1-\alpha_1^2)]$.

If the BS wants the MSS to feed back power weighting for up to $N_s=4$ streams, one 6-bit CQI and one 3-bit CQI (or three 3-bit CQI channels) are allocated. The first power weighting is quantized using 4 bits, the second using 3 bits, and the third using 2 bits. A numerical range of $[\sqrt{1/4}, 1]$ is first uniformly divided into $2^4=16$ levels (i.e., with the interval between levels being $(1-\sqrt{1/4})/15$) and the MSS quantizes the squared root of the power weighting of the first stream to the nearest level (denoted as α_1). Then, 3 bits are used to quantize the range of $[\frac{1}{3}(1-\alpha_1^2), \min(\alpha_1^2, 1-\alpha_1^2)]$ for the second stream squared-root power weighting (denotes as α_2). Finally, 2 bits are used to quantize the range of $[\frac{1}{2}(1-\alpha_1^2-\alpha_2^2), \min(\alpha_2^2, 1-\alpha_1^2-\alpha_2^2)]$ for the third stream squared-root power weighting.

[Add a new section 11.7.8.11]

11.7.8.11 Advanced Receiver Capability

This field indicates whether the MSS is advanced receiver capable

| Type | Length | Value | Scope |
|------|--------|---|--------------------|
| 21 | 1 | Bit 0: Successive Interference Receiver Capability Bit 1-7: Reserved | REG-REQ REG-RSP |

----- End of Text Changes for Remedy 2 -----

References

- [1] T. A. Thomas and F. W. Vook, "A Method for Improving the Performance of Successive Cancellation in Mobile Spread MIMO OFDM," *Proc. IEEE VTC-2002/Fall*, Vancouver, Canada, September 2002.