## IEEE P802.15 Wireless Personal Area Networks

Project	IEEE P802.15 Working Group for Wireless Personal Area Networks (WPANs)		
Title	CONVERGENCE OF THE INSTANT CHANNEL REPLACEMENT ALGORITHM (ACL + SCO – HV2 LINK)		
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Re:	[]		
Abstract	Convergence of the adaptive frequency hopping algorithm for ACL + SCO HV2 traffic is analyzed in the Bandspeed Adaptive Frequency Hopping – Instant Channel Replacement proposal.		
Purpose	Clarification for TG2 members.		
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#### DEFINITIONS

Transaction	a pair of Tx – Rx slots
Ν	total number of hop channels
N <sub>G</sub>	number of 'Good' channels
N <sub>BK</sub>	number of 'Bad-To-Keep' channels
N <sub>BN</sub>	number of 'Bad-To-Replace' channels

#### 1. PROBABILITIES

1.1. The probability of generating a 'Good' channel

$$P_G = \frac{N_G}{N} \tag{1}$$

1.2. The probability of generating a 'Bad-To-Keep' channel

$$P_{BK} = \frac{N_{BK}}{N} \tag{2}$$

1.2. The probability of generating a 'Bad-To-Replace' channel

$$P_{BN} = \frac{N_{BN}}{N} \tag{3}$$

1.3. The probability of a 'Bad-To-Replace' channel is the sum of the probability that a 'Bad-To-Replace' channel will be replaced by 'Good' channel and the probability that 'Bad-To-Replace' channel will be replaced by 'Bad-To-Keep' channel.

$$P_{BN} = P_{BN \to G} + P_{BN \to BK} \tag{4}$$

1.4. The probability that 'Bad-To-Replace' channel will be replaced by 'Good' channel (see App. 1)

$$P_{BN \to G} = \frac{1}{1+K} \cdot P_{BN} \tag{5}$$

1.5. The probability that 'Bad-To-Replace' channel will be replaced by 'Bad-To-Keep' channel (see App. 1)

$$P_{BN \to BK} = \frac{K}{1+K} \cdot P_{BN} , \qquad (6)$$
$$K = \frac{N_{BK}}{N_G}$$

1.6. Total probabilities of appearance of 'Good' and 'Bad-To-Keep' channels

$$P_{\Sigma G} = P_G + P_{BN \to G} \tag{7}$$

$$P_{\sum BK} = P_{BK} + P_{BN \to BK} \tag{8}$$

1.7. Transaction probabilities: 'GG', 'GBK', 'BKG', 'BKBK'

$$P_{GG} = P_{\Sigma G} \cdot P_{\Sigma G} \tag{9}$$

$$P_{GBK} = P_{\Sigma G} \cdot P_{\Sigma BK} \tag{10}$$

$$P_{BKG} = P_{\Sigma BK} \cdot P_{\Sigma G} \tag{11}$$

$$P_{BKBK} = P_{\sum BK} \cdot P_{\sum BK}$$
(12)

Note: 
$$P_{GBK} = P_{BKG}$$

1.8. Probabilities in ACL + SCO (HV2) link.

In an ACL + SCO (HV2) link, half of the Tx - Rx pairs (transactions) are allocated for voice and another half of transactions are allocated for data.

HV2 transaction probabilities:

ACL transaction probabilities:

$P_{GG_{HV2}} = 0.5 P_{GG};$	$P_{GG_{-}ACL} = 0.5 P_{GG};$
$P_{GBK\_HV2} = 0.5 P_{GBK};$	$P_{GBK\_ACL} = 0.5 P_{GBK};$
$P_{BKG_HV2} = 0.5 P_{BKG};$	$P_{BKG\_ACL} = 0.5P_{BKG};$
$P_{BKBK \_ HV2} = 0.5 P_{BKBK}$	$P_{BKBK\_ACL} = 0.5 P_{BKBK}$

1.9. The probability of appearance of 'Bad-To-Keep' channels  $(B_K \rightarrow G \text{ replacements})$  in HV2 slots

$$P_{R_{-}HV2} = P_{GBK_{-}HV2} + P_{BKG_{-}HV2} + P_{BKBK_{-}HV2}$$
  
= 0.5(P\_{GBK} + P\_{BKG} + P\_{BKBK}) (13)

1.10. The probability of appearance of 'Good' channels ( $G \rightarrow B_K$  replacements) in ACL slots

$$P_{A_{A}CL} = P_{GBK_{A}CL} + P_{BKG_{A}CL} + P_{GG_{A}CL}$$

$$= 0.5(P_{GBK} + P_{BKG} + P_{GG})$$
(14)

## 2. CONVERGENCE WHEN $N_G \ge N_{BK}$

For the ICR algorithm to converge (the GUD does not grow indefinitely large) we require the probability of appearance of 'Bad-To-Keep' channels in HV2 slots to be smaller than or equal to the probability of appearance of 'Good' channels in ACL slots

Submission

$$P_{R_{-}HV2} \le P_{A_{-}ACL} \tag{15}$$

or, equivalently, (by substitution of (13) and (14) into (15))

$$P_{GG} \ge P_{BKBK} \tag{16}$$

Condition (16) is satisfied if

$$N_G \ge N_{BK} \tag{17}$$

(See App. 2 for proof).

## 3. CONVERGENCE WHEN $N_G \leq N_{BK}$

To provide convergence when  $N_G \leq N_{BK}$ , do the following first in the low priority timeslots:

- Replace 'Good Good' channel pair to 'Bad Bad' channel pair, to save 2 good channel usage, *i.e.*, decrease Good Channel Usage Debt Counter (GUD) by 2.
- Replace 'Good Bad' channel pair to 'Bad Bad' channel pair, to save 1 good channel usage, *i.e.*, decrease GUD by 1.

In the high priority timeslots, do the following:

- Keep 'Good Good' channel pair untouched
- Replace 'Good Bad' pair to 'Good Good' channel pair as usual
- For 'Bad Good' channel pair, if GUD < -1, then replace to 'Good Good' channel pair, and increment GUD by 1. If GUD > -1, replace to 'Bad Bad' as usual.
- For 'Bad Bad' channel pair, if GUD < -1, then replace to 'Good Good' channel pair, and increment GUD by 2. If GUD > -1, keep it untouched as usual.

By this way, GUD is always converged towards to zero.

#### APPENDIX 1

# $P_{BN \rightarrow G}$ AND $P_{BN \rightarrow BK}$ PROBABILITIES

 $P_{BN \to G}$  and  $P_{BN \to BK}$  may be obtained from the following system of equations

$$\begin{cases} P_{BN} = P_{BN \to G} + P_{BN \to BK} \\ P_{BN \to BK} = K \cdot P_{BN \to G} \end{cases}$$
(1A)

Where K is the proportionality coefficient

$$K = \frac{N_{BK}}{N_G}$$

Solving the system (1A) subject to  $P_{BN \rightarrow G}$  we get

$$P_{BN \to G} = \frac{1}{1+K} \cdot P_{BN}$$

By analogy

$$P_{BN \to BK} = \frac{K}{1+K} \cdot P_{BN}$$

## APPENDIX 2

# THE PROOF

$$P_{GG} \ge P_{BKBK}$$
 is true if  $N_G \ge N_{BK}$ 

$$P_{GG} = P_{\Sigma G} \cdot P_{\Sigma G} = (P_G + P_{BN \to G})^2$$

$$P_{BKBK} = P_{\sum BK} \cdot P_{\sum BK} = (P_{BK} + P_{BN \to BK})^2$$

So the condition is

$$(P_G + P_{BN \to G})^2 \ge (P_{BK} + P_{BN \to BK})^2$$

$$P_G^2 + 2P_G P_{BN \to G} + P_{BN \to G}^2 \ge P_{BK}^2 + 2P_G P_{BN \to BK} + P_{BN \to BK}^2;$$

$$\frac{N_G^2}{N^2} + \frac{2N_G^2 N_{BN}}{N^2 (N_G + N_{BK})} + \frac{N_{BN}^2 N_G^2}{N^2 (N_G + N_{BK})^2} \ge \frac{N_{BK}^2}{N^2} + \frac{2N_{BK}^2 N_{BN}}{N^2 (N_G + N_{BK})} + \frac{N_{BN}^2 N_{BK}^2}{N^2 (N_G + N_{BK})^2};$$

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$$N_{G}^{2}\left[\frac{1}{N^{2}} + \frac{2N_{BN}}{N^{2}(N_{G} + N_{BK})} + \frac{N_{BN}^{2}}{N^{2}(N_{G} + N_{BK})^{2}}\right] \geq N_{BK}^{2}\left[\frac{1}{N^{2}} + \frac{2N_{BN}}{N^{2}(N_{G} + N_{BK})} + \frac{N_{BN}^{2}}{N^{2}(N_{G} + N_{BK})^{2}}\right];$$

$$N_G^2 \ge N_{BK}^2$$

 $\boldsymbol{N}_{G}$  ,  $\boldsymbol{N}_{BK}$  are always positive, so

$$N_G \geq N_{BK}$$